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Lecture 26: More on Algorithms for Sorting

CS 1110

Introduction to Computing Using Python

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Announcements

- Discussion sections this week
 - First 10 minutes dedicated to getting your started on A6
 - Remaining time is office hour for your A6/Prelim 2 questions
- Final Exam on May 21st 1:30-4pm. Your assigned exam session (in-person or online) is shown in CMS. *Submit a “regrade request” in CMS by May 12 if you have a legitimate reason for requesting a change. If you have an exceptional circumstance for switching from in-person to online, you must upload to CMS your college’s approval of your modality change.*

More Announcements

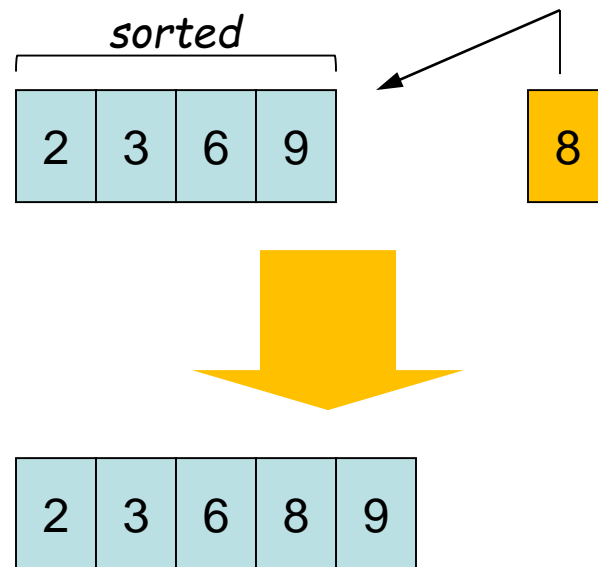
- **A6** due on Friday
 - Remember academic integrity
- Expected release dates of solutions and feedback
 - **A5** solutions: Wed May 12
 - **A4** grades and feedback: Thurs May 13
 - **A6** solutions: Tues May 18
 - **A5** grades and feedback: Thurs May 20
 - **Final exam** grades and feedback: Tues May 25
 - **A6** grades and feedback: Fri May 28

Algorithms for Sorting

- Well known algorithms
 - focus on reviewing programming constructs (**while** loop) and analysis
 - will not use built-in methods such as **sort**, **index**, **insert**, etc.
- Today we'll discuss **merge sort** and compare it to **insertion sort**, which we discussed last lecture
- More on the topic in next course, CS 2110!

The Insertion Process of Insertion Sort

- Given a sorted list x , insert a number y such that the result is sorted
- Sorted: arranged in ascending (small to big) order



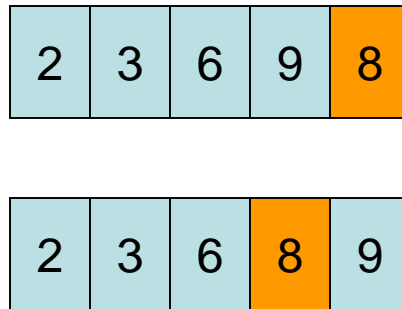
We'll call this process a “**push down**,” as in push a value down until it is in its sorted position

Algorithm Complexity

- Count the number of comparisons needed
- In the worst case, need i comparisons to push down an element in a sorted segment with i elements.

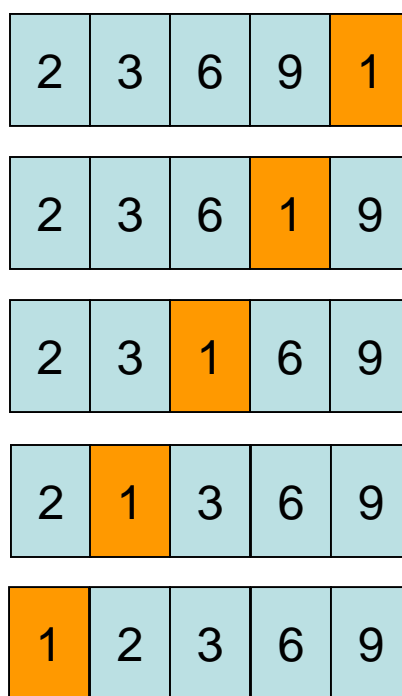
How much work is a push down?

push down
a "big"
value



This push down takes
2 comparisons

push down
a "small"
value



This push down takes
4 comparisons.
Worst case scenario:
 n comparisons
needed to push down
into a length n sorted
segment.

Algorithm Complexity (Q)

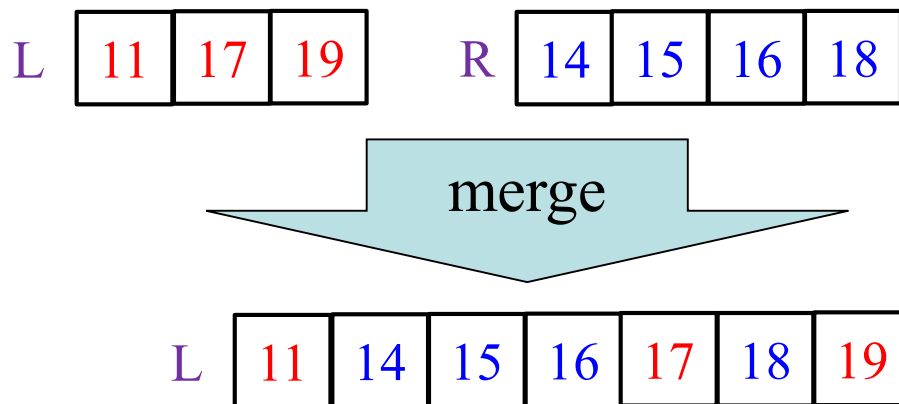
```
def swap(b, h, k):  
    :  
  
def push_down(b, k):  
    while k > 0 and b[k-1] > b[k]:  
        swap(b, k-1, k)  
        k = k-1  
  
def insertion_sort(b):  
    for i in range(1, len(b)):  
        push_down(b, i)
```

Count (approximately) the number of **comparisons** needed to sort a list of length n

- A. ~ 1 comparison
- B. $\sim n$ comparisons
- C. $\sim n^2$ comparisons
- D. $\sim n^3$ comparisons
- E. I don't know

Which algorithm does Python's sort use?

- Recursive algorithm that runs much faster than insertion sort for the same size list (when the size is big)!
- A variant of an algorithm called “merge sort”
- Based on the idea that sorting is hard, but “merging” two *already sorted* lists is easy.



Merge sort: Motivation

Since merging is easier than sorting, if I have two helpers, I'd...

- Give each helper half the array to sort
- Then I get back their sorted subarrays and **merge** them.

What if those two helpers each had two sub-helpers?

And the sub-helpers each had two sub-sub-helpers? And...

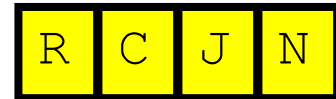
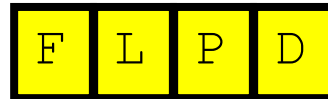
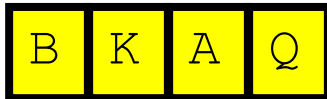
Subdivide the sorting task

H	E	M	G	B	K	A	Q	F	L	P	D	R	C	J	N
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

H	E	M	G	B	K	A	Q
---	---	---	---	---	---	---	---

F	L	P	D	R	C	J	N
---	---	---	---	---	---	---	---

Subdivide again



And again



H E M G

B K A Q

F L P D

R C J N

H E

M G

B K

A Q

F L

P D

R C

J N

And one last time

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H	E
---	---

M	G
---	---

B	K
---	---

A	Q
---	---

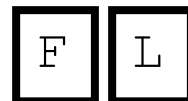
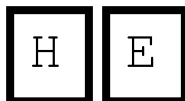
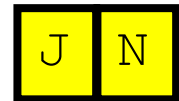
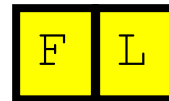
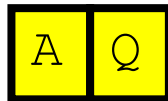
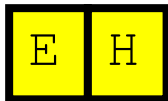
F	L
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P	D
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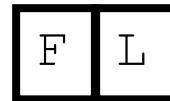
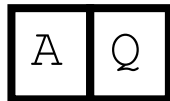
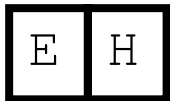
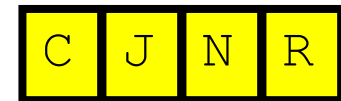
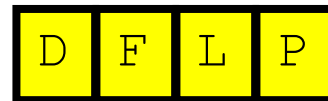
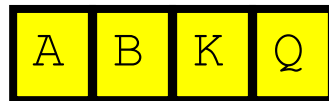
R	C
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J	N
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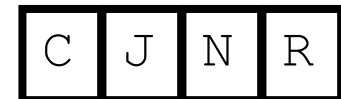
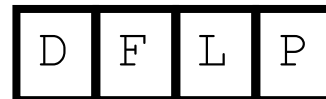
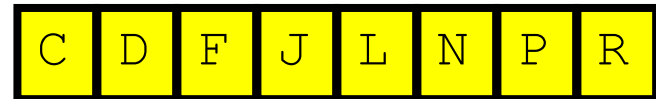
Now merge



And merge again



And again



And one last time

A	B	C	D	E	F	G	H	J	K	L	M	N	P	Q	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

A	B	E	G	H	K	M	Q
---	---	---	---	---	---	---	---

C	D	F	J	L	N	P	R
---	---	---	---	---	---	---	---

Done!

A	B	C	D	E	F	G	H	J	K	L	M	N	P	Q	R
---	---	---	---	---	---	---	---	---	---	---	---	---	---	---	---

```
def mergeSort(li):  
    """Sort list li using Merge Sort"""  
    if len(li) > 1:  
        # Divide into two parts  
        mid= len(li)//2  
        left= li[:mid]  
        right= li[mid:]  
  
        # Recursive calls  
        mergeSort(left)  
        mergeSort(right)  
  
        # Merge left & right back to li  
    ...
```

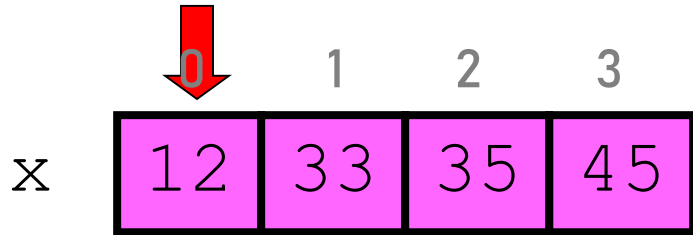
The central sub-problem is the merging of two sorted lists into one single sorted list

12	33	35	45
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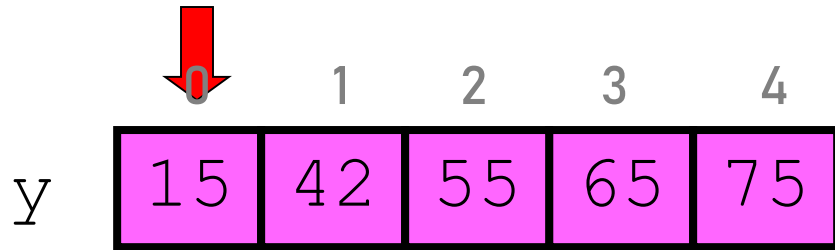
15	42	55	65	75
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12	15	33	35	42	45	55	65	75
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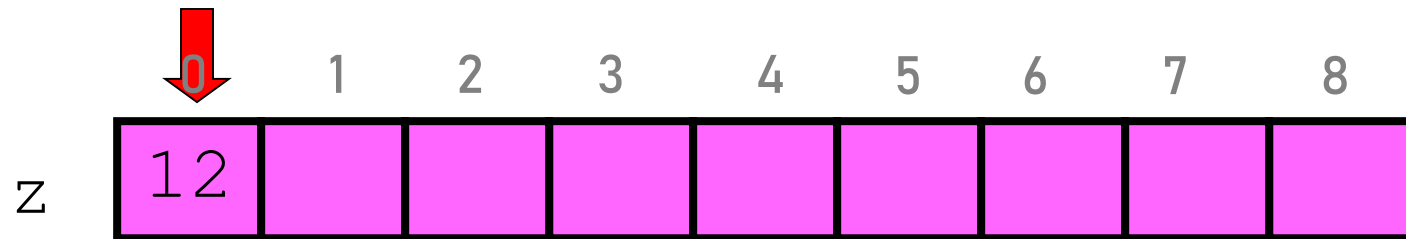
Merge



ix 0



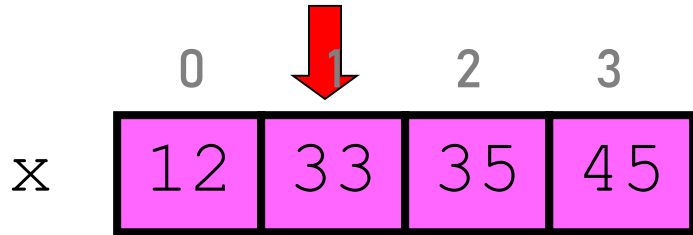
iy 0



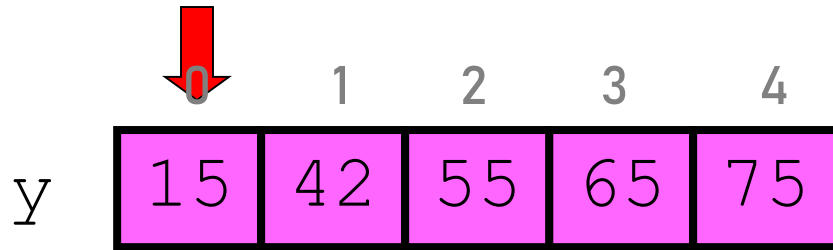
iz 0

$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ YES

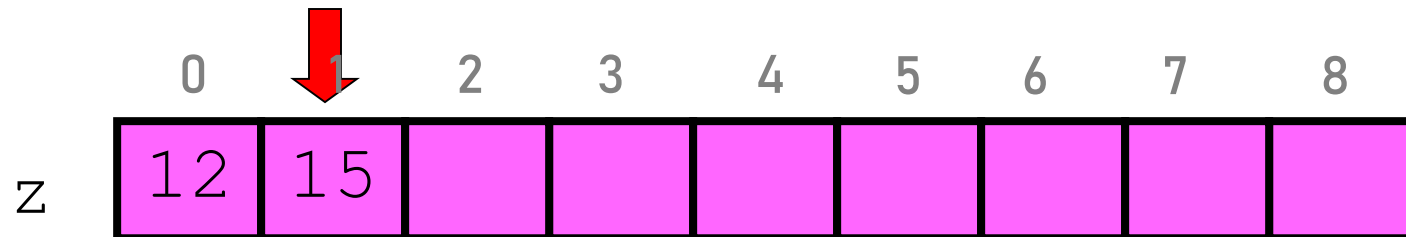
Merge



ix 1



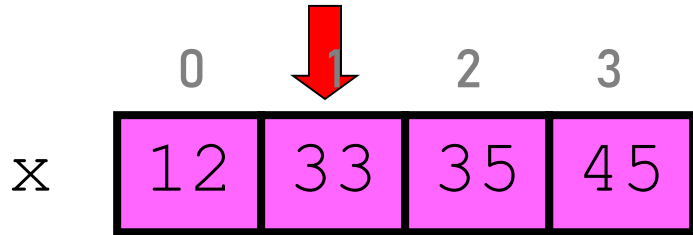
iy 0



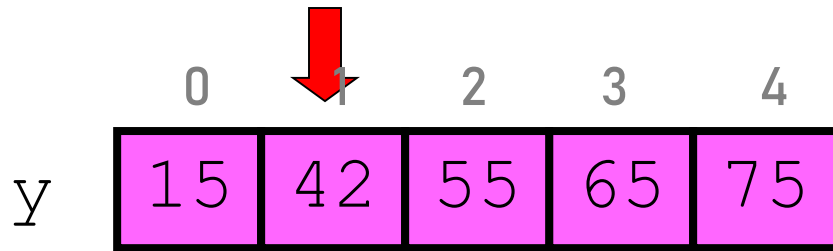
iz 1

$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ NO

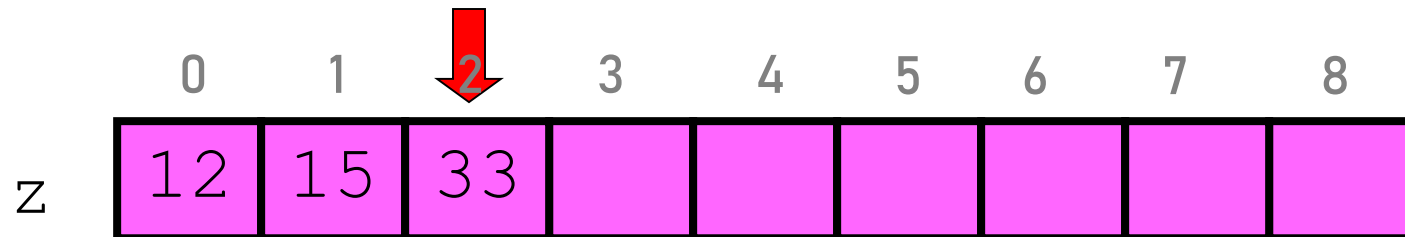
Merge



ix 1



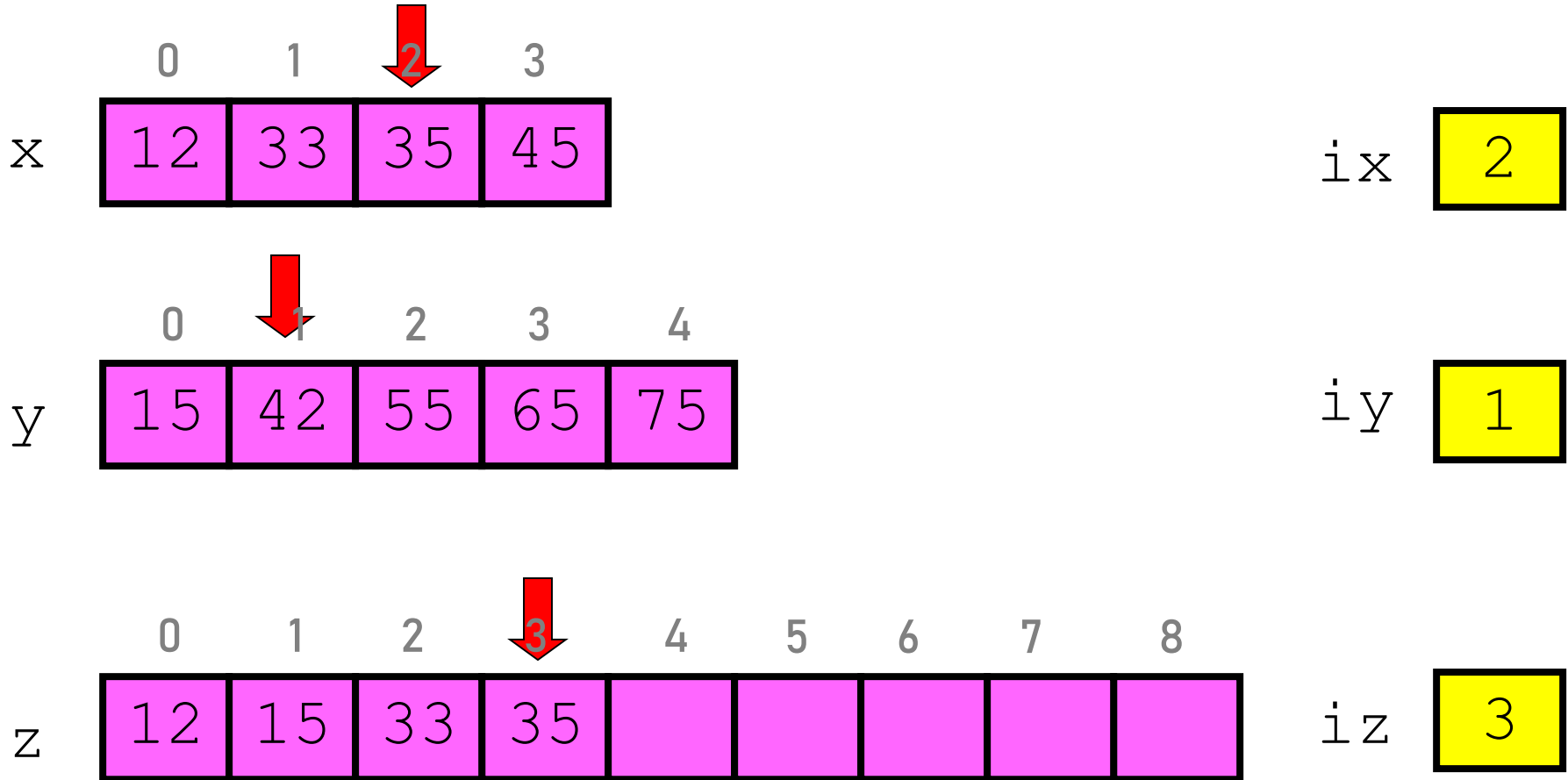
iy 1



iz 2

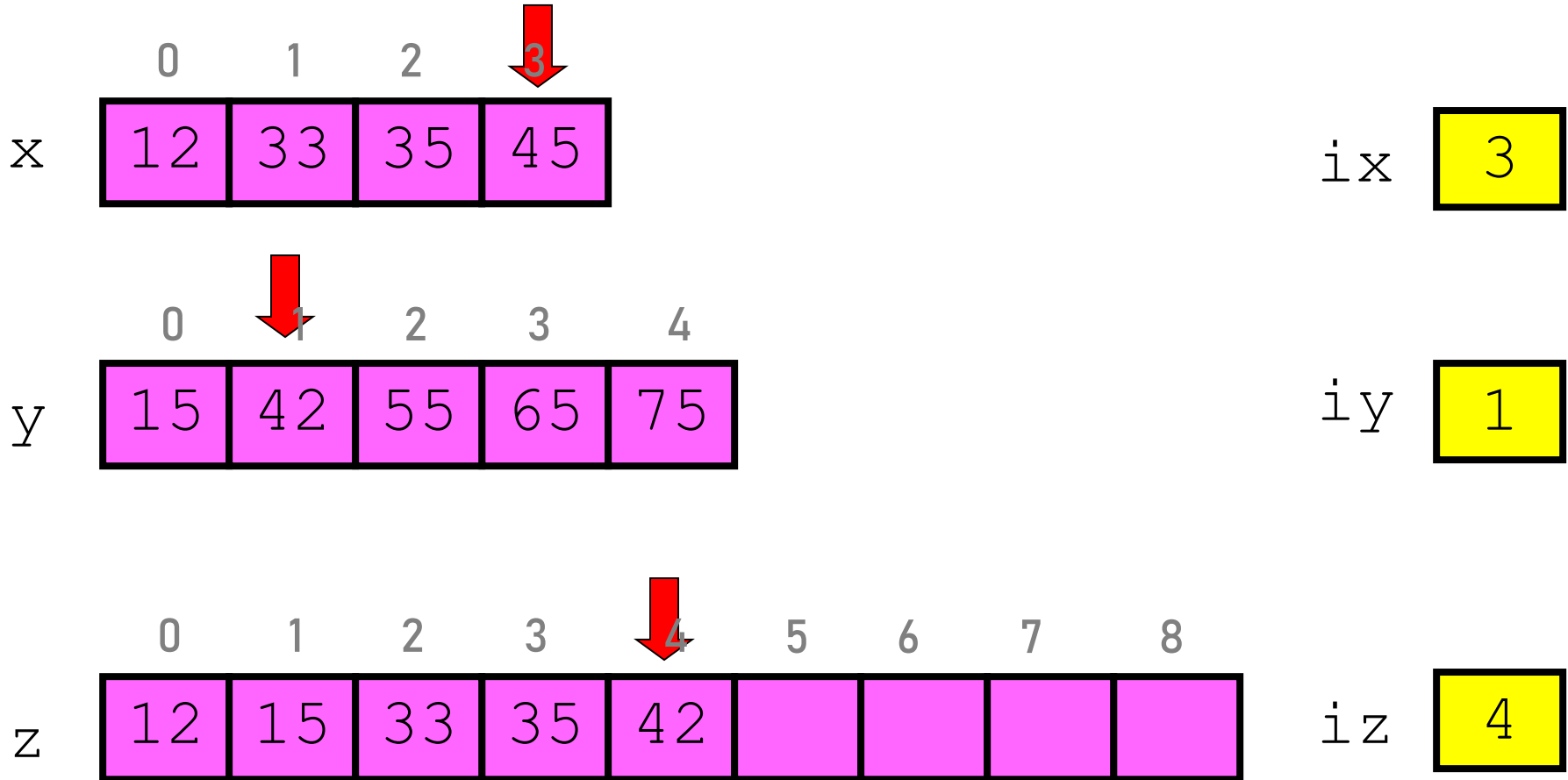
$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ YES

Merge



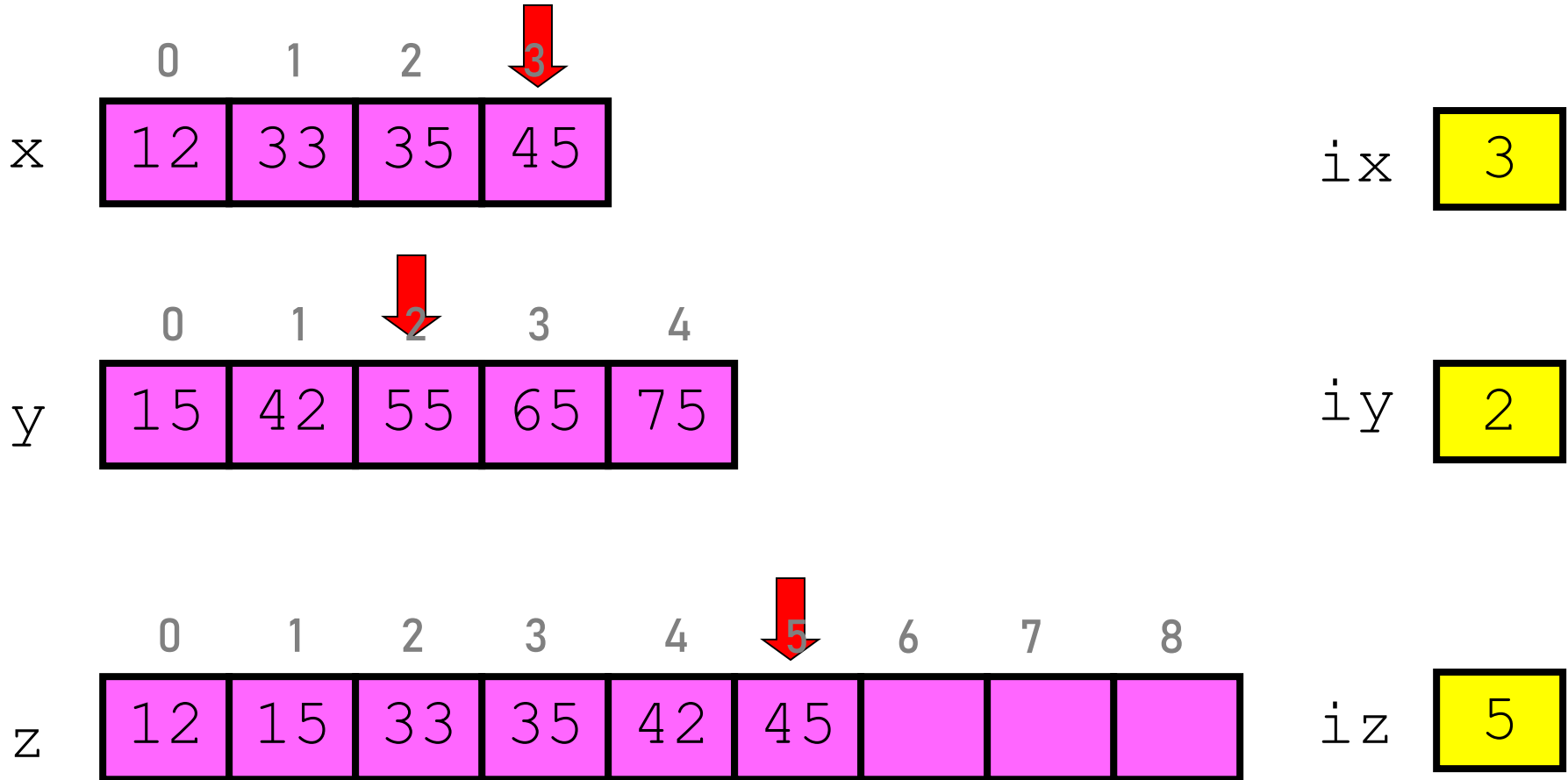
$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ YES

Merge



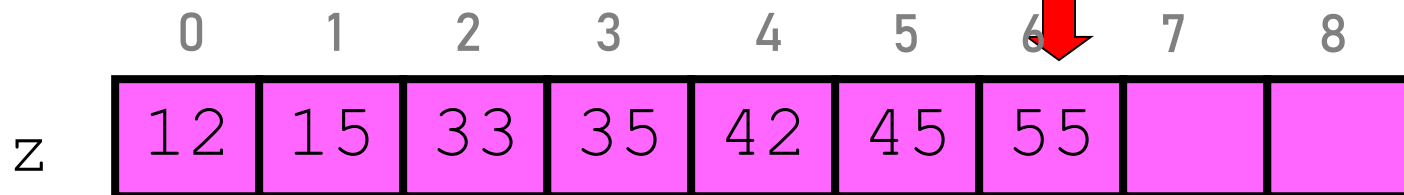
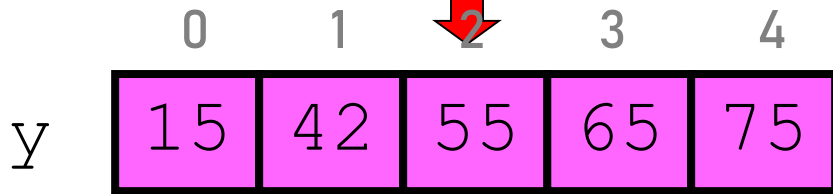
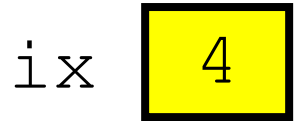
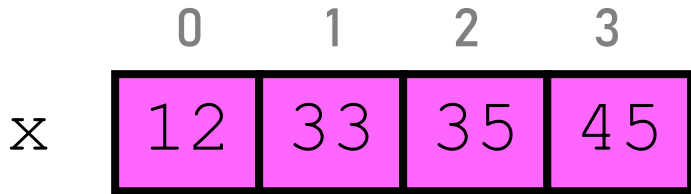
$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ NO

Merge



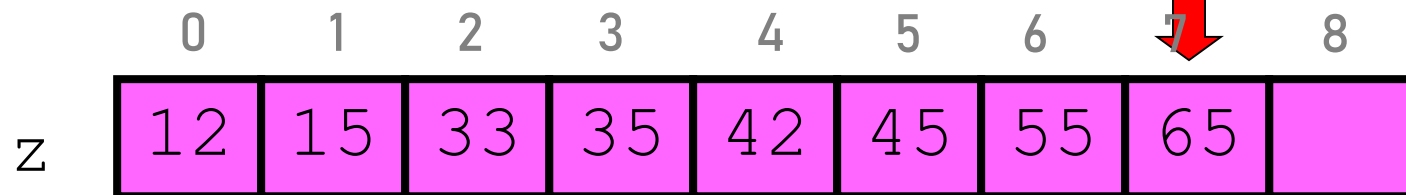
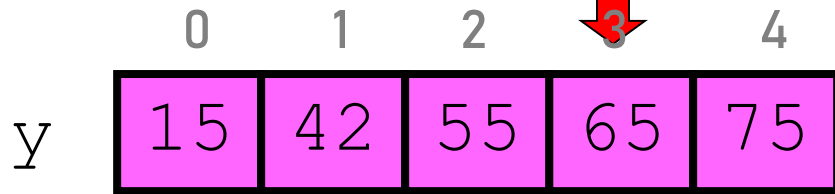
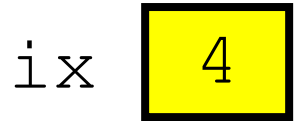
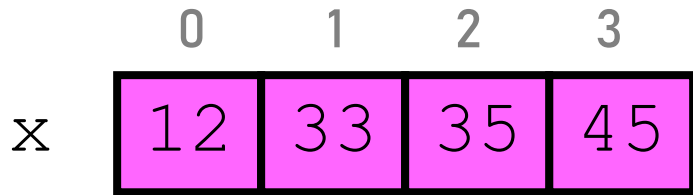
$ix < 4$ and $iy < 5 \rightarrow x(ix) \leq y(iy)$ YES

Merge



ix at 4 → take y(iy)

Merge



$i_y < 5 \rightarrow \text{take } y(i_y)$

Merge



	0	1	2	3
x	12	33	35	45

ix 4

	0	1	2	3	4
y	15	42	55	65	75

iy 4

	0	1	2	3	4	5	6	7	8
z	12	15	33	35	42	45	55	65	75

iz 8

$iy < 5 \rightarrow \text{take } y(iy)$

Merge



	0	1	2	3
x	12	33	35	45

ix	4
----	---



	0	1	2	3	4
y	15	42	55	65	75

iy	5
----	---



	0	1	2	3	4	5	6	7	8
z	12	15	33	35	42	45	55	65	75

iz	9
----	---

```
# Given lists x and y and list z, which has
# the combined length of x and y...
nx = len(x); ny = len(y)

ix = 0; iy = 0; iz = 0;
while ix<nx and iy<ny
    if x[ix] <= y[iy]:
        z[iz]= x[ix]; ix=ix+1
    else:
        z[iz]= y[iy]; iy=iy+1
    iz=iz+1

while ix<nx # copy any remaining x-values
    z[iz]= x[ix]; ix=ix+1; iz=iz+1

while iy<ny # copy any remaining y-values
    z[iz]= y[iy]; iy=iy+1; iz=iz+1
```


How do merge sort and insertion sort compare?

- Insertion sort: (worst case) makes i comparisons to insert an element in a sorted array of i elements. For an array of length n :

_____ for big n

- Merge sort: _____

```
def mergeSort(li):  
    """Sort list li using Merge Sort"""  
    if len(li) > 1:  
        # Divide into two parts  
        mid= len(li)/2  
        left= li[:mid]  
        right= li[mid:]  
  
        # Recursive calls  
        mergeSort(left)  
        mergeSort(right)  
  
        # Merge left & right back to li  
    ...
```



All the comparisons between list elements are done during merge

```
# Given lists x and y and list z, which has  
# the combined length of x and y...
```

```
nx = len(x); ny = len(y)
```

```
ix = 0; iy = 0; iz = 0;
```

```
while ix<nx and iy<ny
```

```
    if x[ix] <= y[iy]:
```

```
        z[iz]= x[ix]; ix=ix+1
```

```
    else:
```

```
        z[iz]= y[iy]; iy=iy+1
```

```
    iz=iz+1
```

```
while ix<nx # copy any remaining x-values
```

```
    z[iz]= x[ix]; ix=ix+1; iz=iz+1
```

```
while iy<ny # copy any remaining y-values
```

```
    z[iz]= y[iy]; iy=iy+1; iz=iz+1
```

Merge – best case scenario

x

2	3	5	14
---	---	---	----

y

15	42	55	65
----	----	----	----

z

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Merge – worst case scenario

x

12	23	45	64
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y

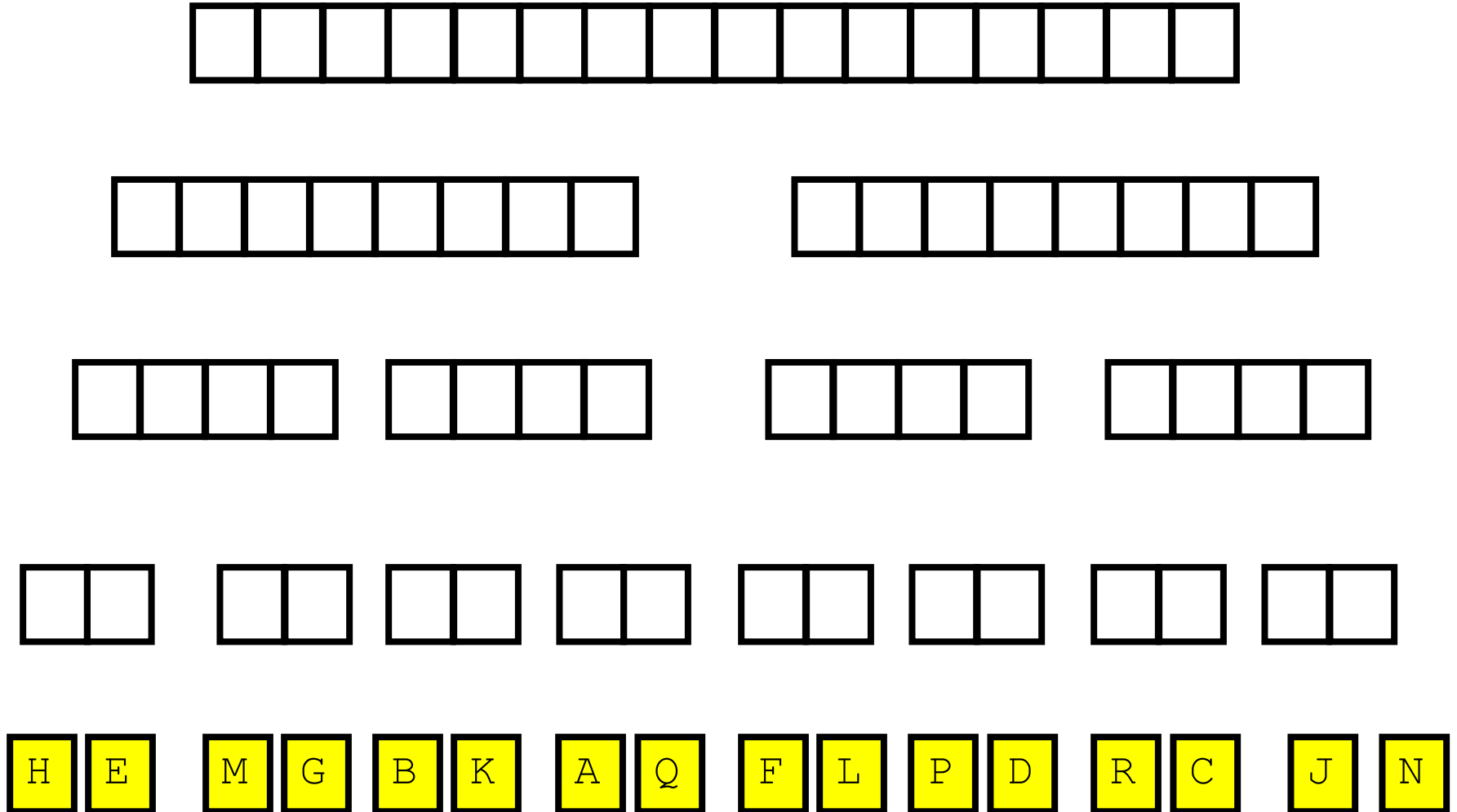
15	42	55	65
----	----	----	----

z

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Need to do $n-1$ comparisons where n is total number of elements merged

Merge sort: about $\log_2(n)$ “levels”;
about n comparisons each level




How do merge sort and insertion sort compare?

- Insertion sort: (worst case) makes i comparisons to insert an element in a sorted array of i elements. For an array of length n :

$$1+2+\dots+(n-1) = n(n-1)/2, \text{ say } n^2 \text{ for big } n$$

- Merge sort: $n \cdot \log_2(n)$ comparisons

Order of
magnitude
difference



- Should we always use merge sort then? Python actually uses a variant that combines merge sort and insertion sort!