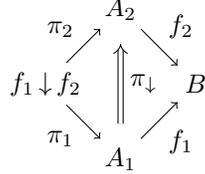


Weighted Limits and Colimits

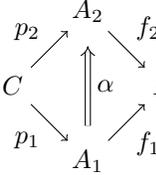
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Definition (Comma Object). Given 1-cells $A_1 \xrightarrow{f_1} B \xleftarrow{f_2} A_2$ of a 2-category, a comma object from f_1 to f_2 is a 0-cell, typically denoted $f_1 \downarrow f_2$, along with 1- and 2-cells as in the following diagram



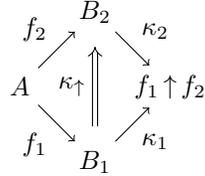
that is *universal* in the sense that given any other diagram $C \begin{array}{ccc} & A_2 & \\ & \uparrow \alpha & \\ C & & B \end{array}$ there exists a unique 1-cell $\langle \alpha \rangle : C \rightarrow f_1 \downarrow f_2$



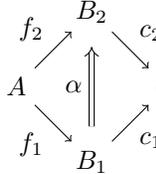
such that $\langle \alpha \rangle * \pi_{\downarrow}$ equals α (and $\langle \alpha \rangle ; \pi_1$ equals p_1 and $\langle \alpha \rangle ; \pi_2$ equals p_2).

Example. Comma categories are the comma objects of **Cat**.

Definition (Cocomma Object). Given 1-cells $B_1 \xleftarrow{f_1} A \xrightarrow{f_2} B_2$ of a 2-category, a cocomma object from f_1 to f_2 is a 0-cell, typically denoted $f_1 \uparrow f_2$, along with 1- and 2-cells as in the following diagram



that is *universal* in the sense that given any other diagram $A \begin{array}{ccc} & B_2 & \\ & \uparrow \alpha & \\ A & & C \end{array}$ there exists a unique 1-cell $[\alpha] : f_1 \uparrow f_2 \rightarrow C$



such that $\kappa_{\uparrow} * [\alpha]$ equals α (and $\kappa_1 ; [\alpha]$ equals c_1 and $\kappa_2 ; [\alpha]$ equals c_2).

Example. For the 1-source $1 \xleftarrow{!} A \xrightarrow{id} A$ in **Prost**, the corresponding cocomma object $! \uparrow A$ is the set $\text{Option}(A)$ with none being smaller than $\text{some}(a)$ for all $a \in A$. On the flipside, the cocomma object $A \uparrow !$ is the set $\text{Option}(A)$ with none being larger than $\text{some}(a)$ for all $a \in A$. In both cases, $\text{some}(a)$ is less than $\text{some}(a')$ iff a is less than a' .

Note that $\mathbb{L}(A)$ in **Set** can be defined as the fixpoint $\mu X. 1 + (A \times X)$. In **Prost**, the fixpoints $\mu X. 1 + (A \times X)$, $\mu X. ! \uparrow (A \times X)$, and $\mu X. (A \times X) \uparrow !$ all correspond to lists but with different orderings. In the first, $\ell \leq \ell'$ can only hold if ℓ and ℓ' have the same length, whereas in the second ℓ can be a prefix of ℓ' , and in the third ℓ' can be a prefix of ℓ . In particular, they all agree on lists with the same length, in which case they use componentwise comparison; where they differ is how they handle lists of differing length.

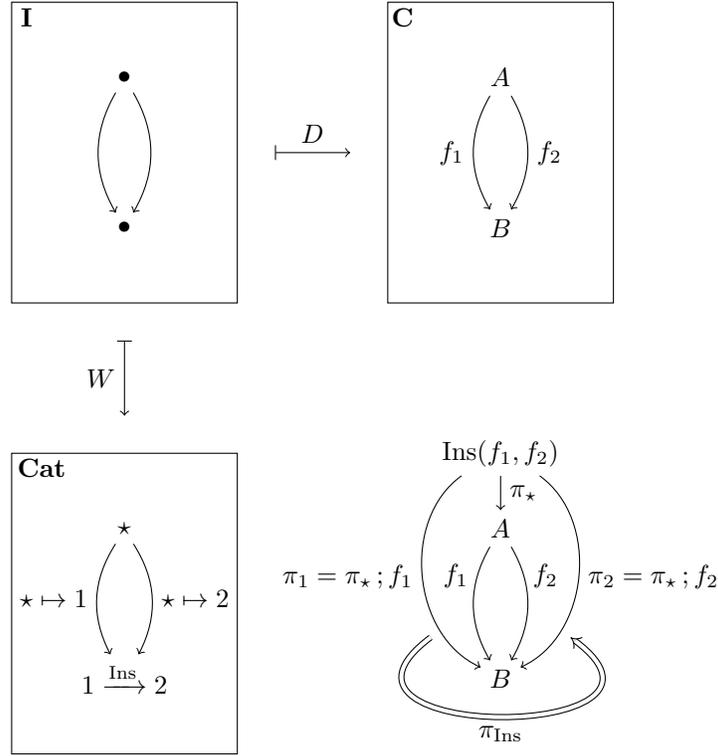
Definition (Inserter Object). Given two 1-cells $f_1, f_2 : A \rightarrow B$ of a 2-category, an inserter from f_1 to f_2 is a 0-cell, typically denoted $\text{Ins}(f_1, f_2)$, along with 1-cell $\pi : \text{Ins}(f_1, f_2) \rightarrow A$ and 2-cell $\pi_{\text{Ins}} : \pi ; f_1 \Rightarrow \pi ; f_2 : \text{Ins}(f_1, f_2) \rightarrow B$ that is *universal*, meaning given any other 0-cell C with 1-cell $f : C \rightarrow A$ and 2-cell $\alpha : f ; f_1 \Rightarrow f ; f_2 : C \rightarrow B$ there exists a unique 1-cell $\langle \alpha \rangle : C \rightarrow \text{Ins}(f_1, f_2)$ such that $\langle \alpha \rangle ; \pi$ equals f and $\langle \alpha \rangle * \pi_{\text{Ins}}$ equals α .

Example. Given an endofunctor $T : \mathbf{C} \rightarrow \mathbf{C}$, the category $\mathbf{Alg}(T)$ is the inserter from T to $\text{Id}_{\mathbf{C}}$, and the category $\mathbf{Coalg}(T)$ is the inserter from $\text{Id}_{\mathbf{C}}$ to T .

Definition (Coinsserter Object). Given two 1-cells $f_1, f_2 : A \rightarrow B$ of a 2-category, a coinsserter from f_1 to f_2 is a 0-cell, $\text{Coins}(f_1, f_2)$, along with 1-cell $\kappa : B \rightarrow \text{Coins}(f_1, f_2)$ and 2-cell $\kappa_{\text{Coins}} : f_1; \kappa \Rightarrow f_2; \kappa : A \rightarrow \text{Coins}(f_1, f_2)$ that is *(co)universal*, meaning given any other 0-cell C with 1-cell $f : B \rightarrow C$ and 2-cell $\alpha : f_1; f \Rightarrow f_2; f : A \rightarrow C$ there exists a unique 1-cell $[\alpha] : \text{Coins}(f_1, f_2) \rightarrow C$ such that $\kappa; [\alpha]$ equals f and $\kappa_{\text{Coins}} * [\alpha]$ equals α .

Definition (Weighted Limit). Let \mathbf{I} be a 2-category conceptually describing a scheme, and let $D : \mathbf{I} \rightarrow \mathbf{C}$ be a 2-functor conceptually describing a diagram of scheme \mathbf{I} in the 2-category \mathbf{C} . Furthermore, let $W : \mathbf{I} \rightarrow \mathbf{Cat}$ be a 2-functor conceptually describing a *weighting* of the diagram. A W -weighted *cone* of the diagram D , denoted $\lim_W D$, is a 0-cell L of \mathbf{C} and a collection of 1-cells $\{\pi_w : L \rightarrow DI\}_{I \in \mathbf{I}, w \in WI}$ and 2-cells $\{\pi_w : \pi_w \Rightarrow \pi_{w'}\}_{I \in \mathbf{I}, w \rightarrow w' \in WI}$ that preserves identities and compositions, meaning $\pi_{id_w} = id_{\pi_w}$ and $\pi_w; \omega' = \pi_w; \pi_{\omega'}$, and is *natural*, meaning for all 1-cells $i : I \rightarrow I' \in \mathbf{I}$ each appropriate 1-cell $\pi_{(Wi)(w)}$ equals $\pi_w; Di$ and for all 2-cells $\iota : i \Rightarrow i' \in \mathbf{I}$ each appropriate 2-cell $\pi_{(W\iota)_w}$ equals $\pi_w * D\iota$. A W -weighted *limit* of a diagram D is a universal W -weighted cone of D .

Example. An inserter is a weighted limit as illustrated below:



Example. A comma object is a weighted limit. The scheme is $\bullet \rightarrow \bullet \leftarrow \bullet$ and the weighting is $1 \hookrightarrow (1 \downarrow 2) \hookleftarrow 2$.

Definition. A weighted colimit is dual to a weighted limit: given a diagram $D : \mathbf{I} \rightarrow \mathbf{C}$ and weighting $W : \mathbf{I}^{\text{op}} \rightarrow \mathbf{Cat}$ (the reason that W is contravariant here is complicated and very meta), a weighted colimit $\text{colim}_W D$ in \mathbf{C} is a weighted limit $\lim_W D^{\text{op}}$ in \mathbf{C}^{op} .

Example. Coinsserter and cocomma objects are weighted colimits of the same weighting but on the opposite scheme as for inserter and comma objects.